

Towards a Real-Time Systems Compiler

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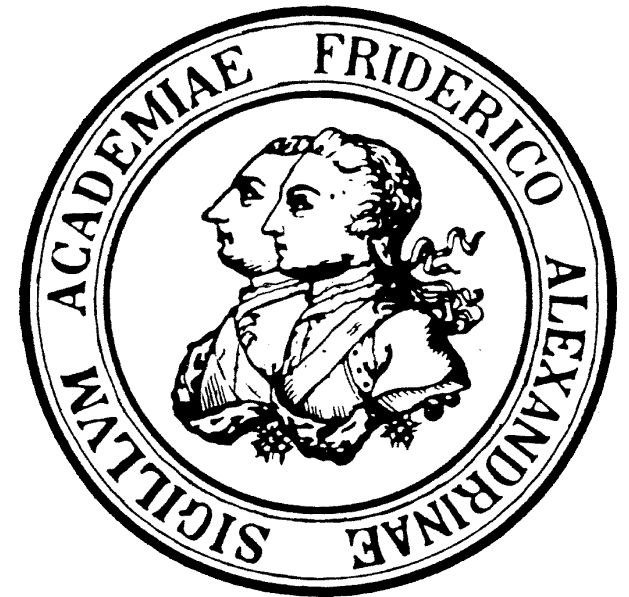
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Recent Cars – Future Cars

■ **event-triggered** communication

▪ mainly CAN



■ **federated** architecture

▪ OSEK/VDX



■ **fail stop** semantics

▪ ABS, ESP



■ **time-triggered** communication

▪ FlexRay



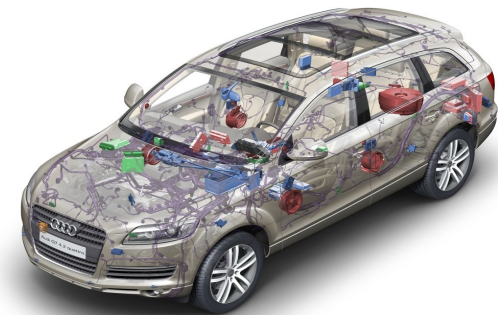
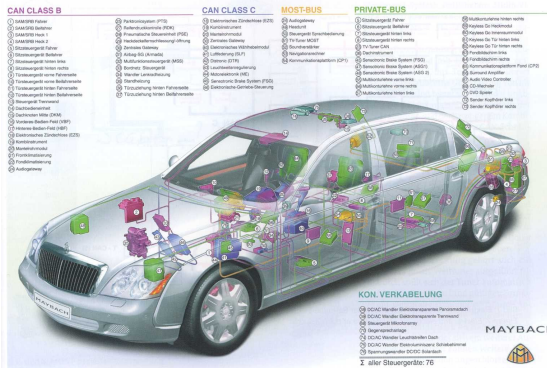
■ **integrated** architecture

▪ AUTOSAR



■ **fail operational** semantics

▪ {Steer,Break}-by-Wire



Consequences

- Migration
 - from an **event-triggered** to a **time-triggered** environment

- Legacy Applications
 - if possible: reuse (e.g. via virtual CAN-networks on top of FlexRay)
 - non-safety-critical subsystems, e.g. comfort applications
 - otherwise: migration/porting
 - safety-critical subsystems, e.g. ABS, ESP

- Porting
 - labour intensive
 - source of errors

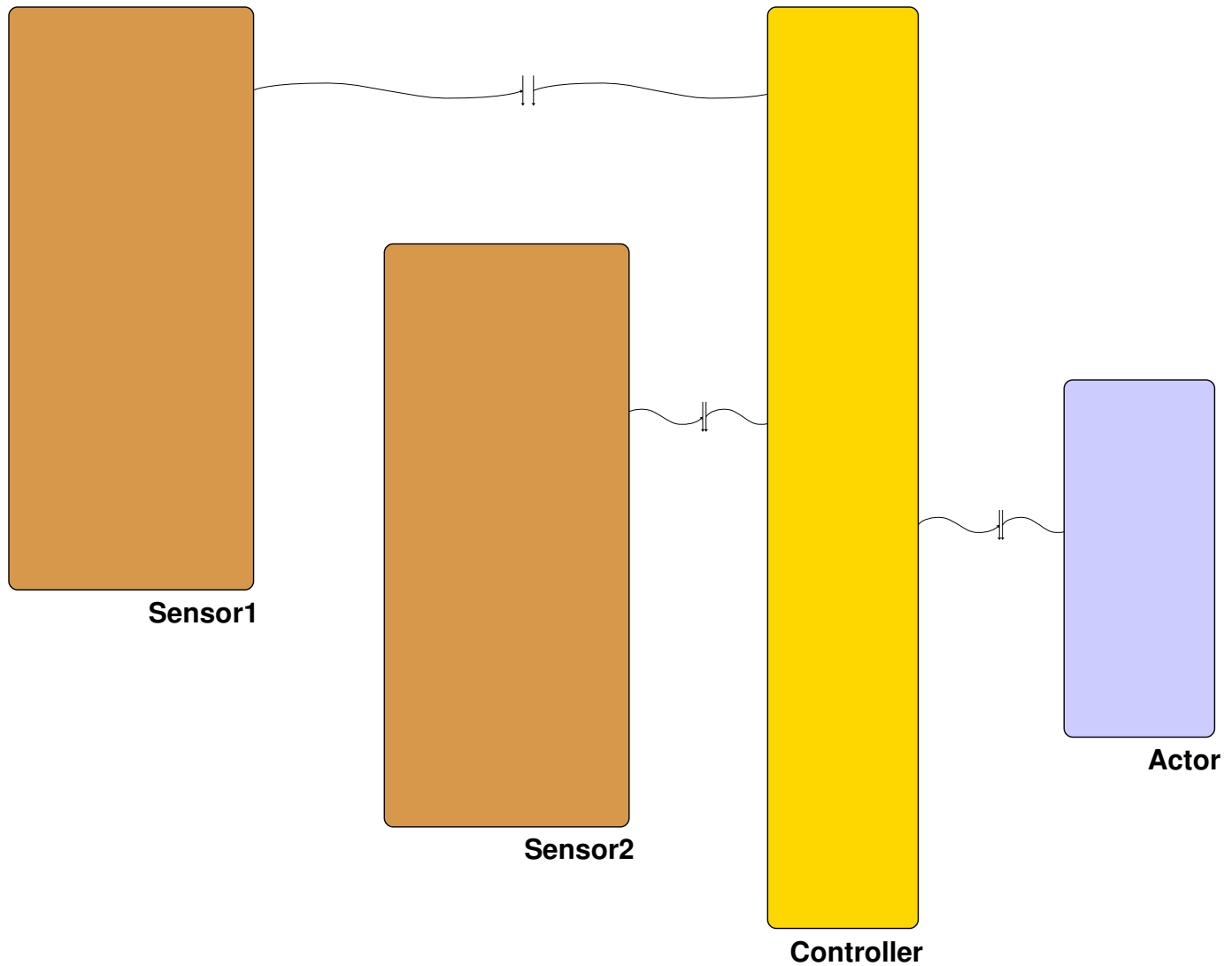


Overview

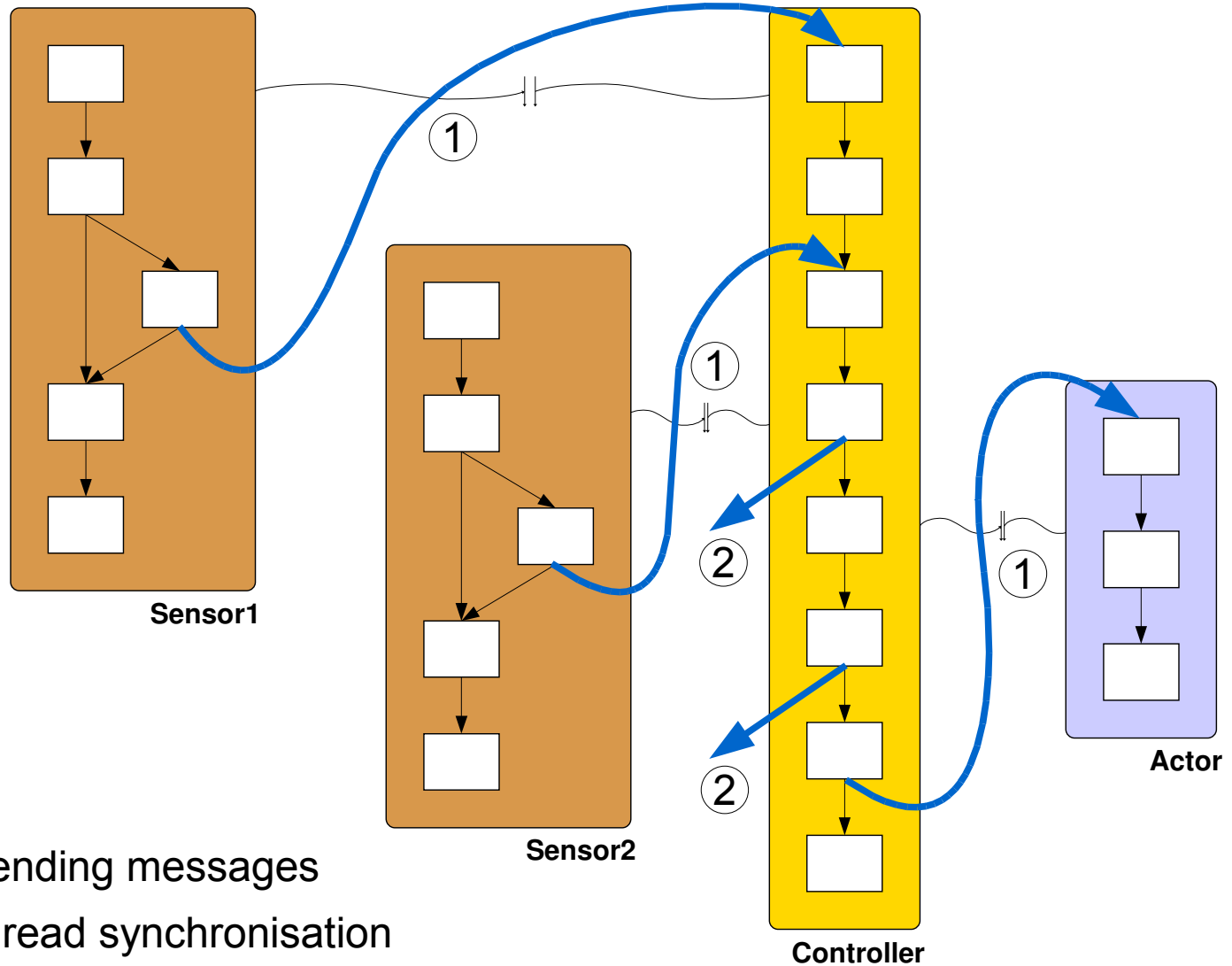
- **Why Migration is a Problem**
- Aiding Migration
- Atomic Basic Blocks
- The Real Time Systems Compiler
- Current Status & Future Work



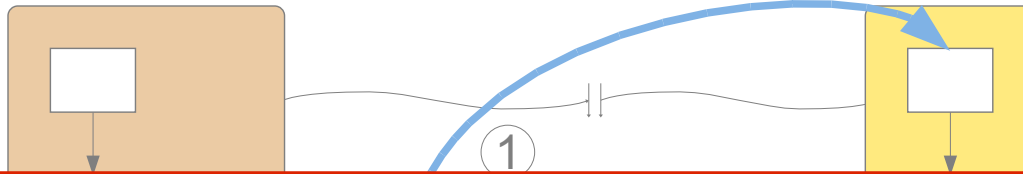
A simple scenario



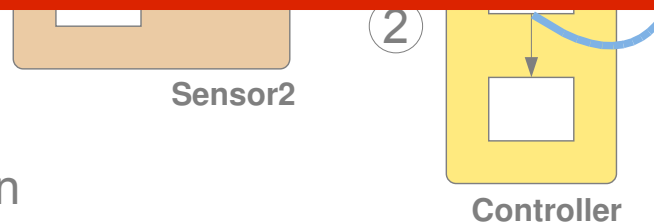
Explicitly Modelled Dependencies



Explicitly Modelled Dependencies



- this is a very synthetic and small example
 - real-world applications contain ≥ 200 Tasks
 - order of magnitude of more dependencies
- computing static schedules is not easy
 - sample another signal at high frequency
 - computations may exceed time slots
 - splitting computations up ... **manually???**



- ① sending messages
- ② thread synchronisation



Overview

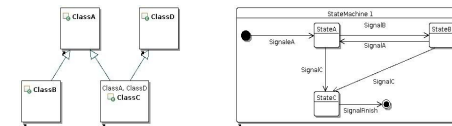
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Lowering Transformations

- automated & generic

High Level of Abstraction



Modelling

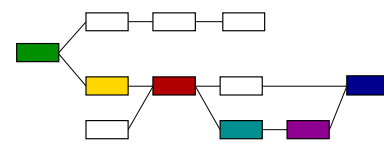
```
int main() {
    // Create a mutex and a semaphore
    pthread_mutex_t mutex = PTHREAD_MUTEX_INITIALIZER;
    pthread_semaphore_t sem = PTHREAD_SEMAPHORE_INITIALIZER;
    sem.value = 5;

    // Create three threads
    pthread_t t1, t2, t3;
    pthread_create(&t1, NULL, thread1, &mutex);
    pthread_create(&t2, NULL, thread2, &sem);
    pthread_create(&t3, NULL, thread3, &sem);

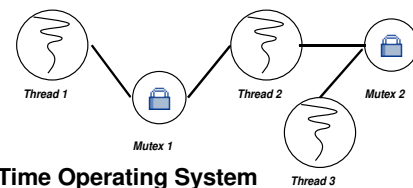
    // Wait for all threads to finish
    pthread_join(t1, NULL);
    pthread_join(t2, NULL);
    pthread_join(t3, NULL);

    return 0;
}
```

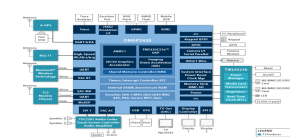
Application



Run-Time System

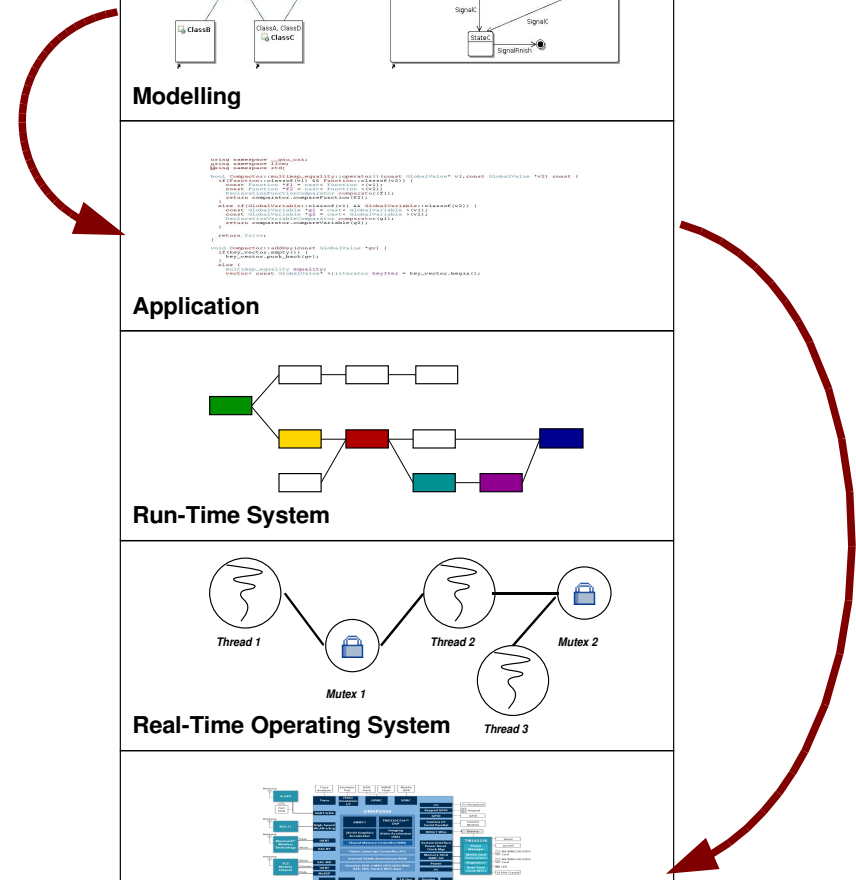


Real-Time Operating System



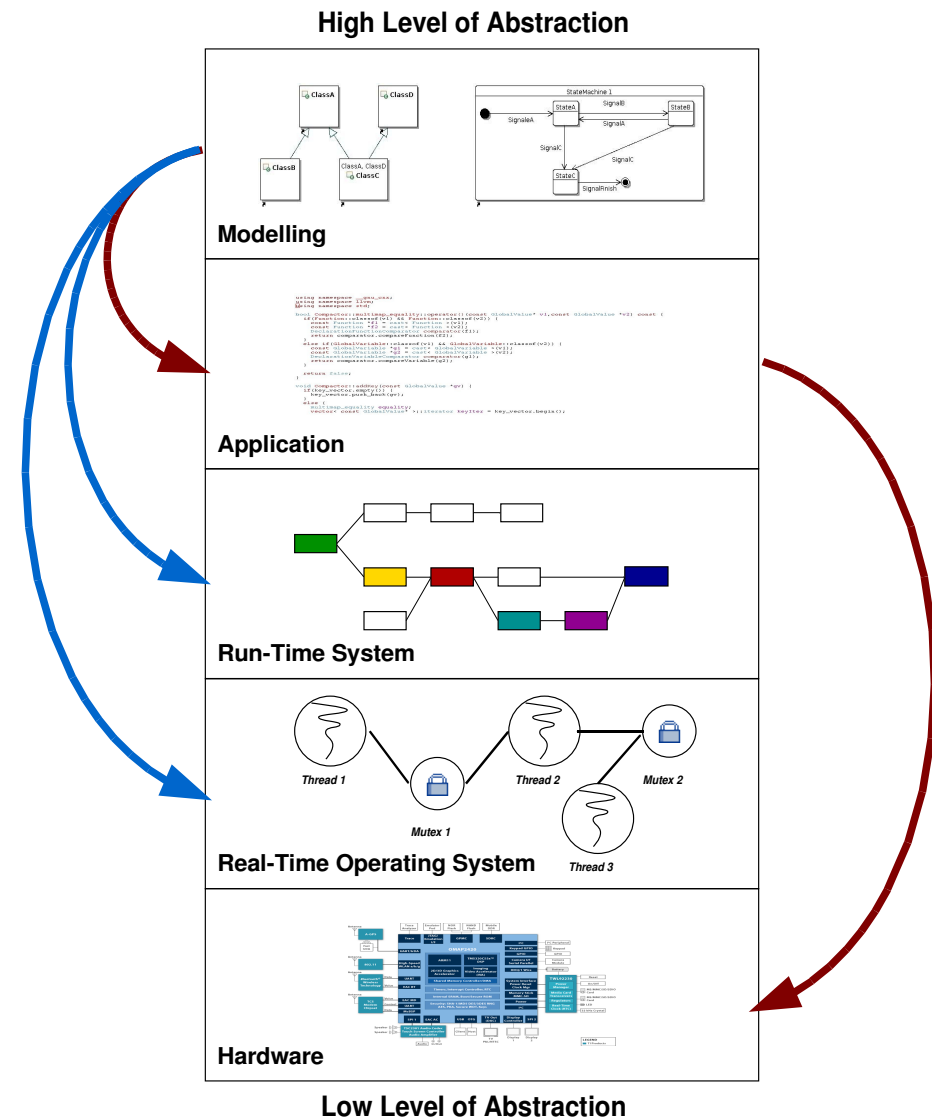
Hardware

Low Level of Abstraction



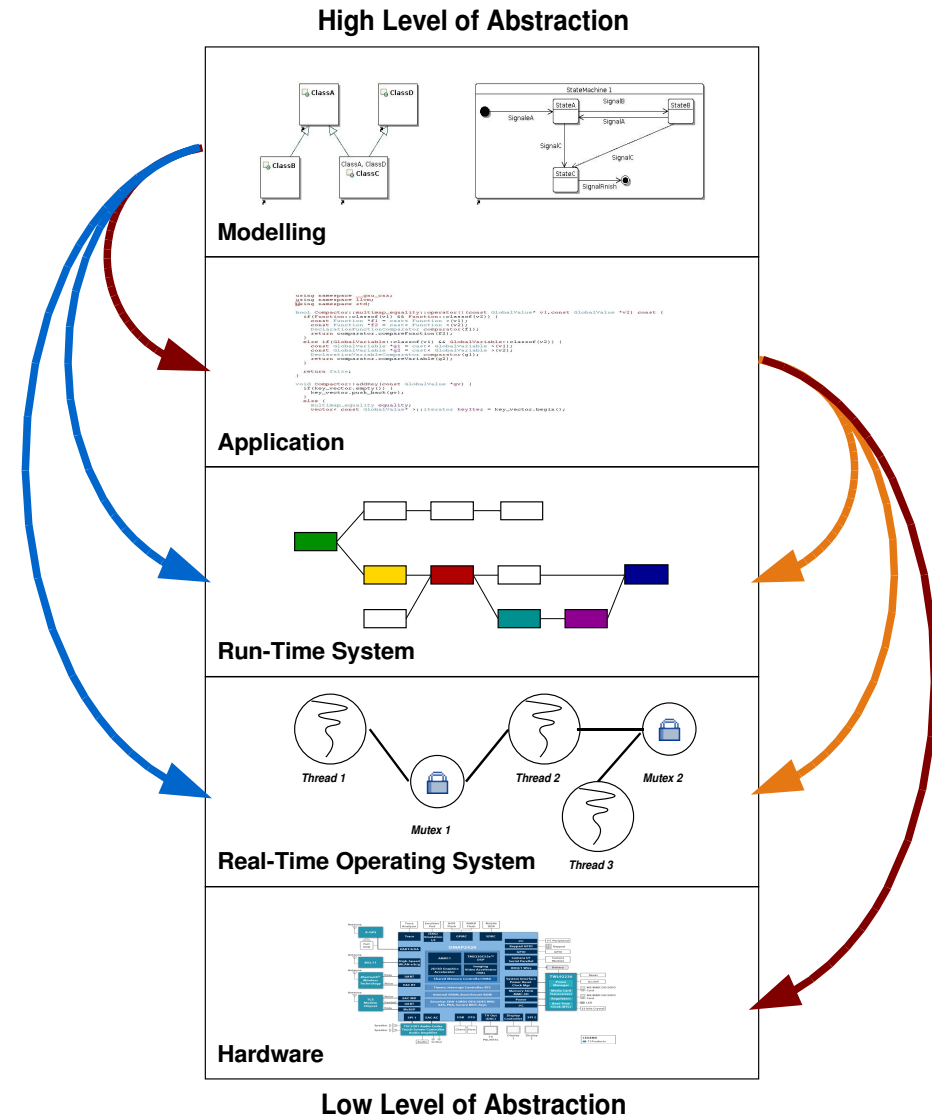
Lowering Transformations

- automated & generic
- automated



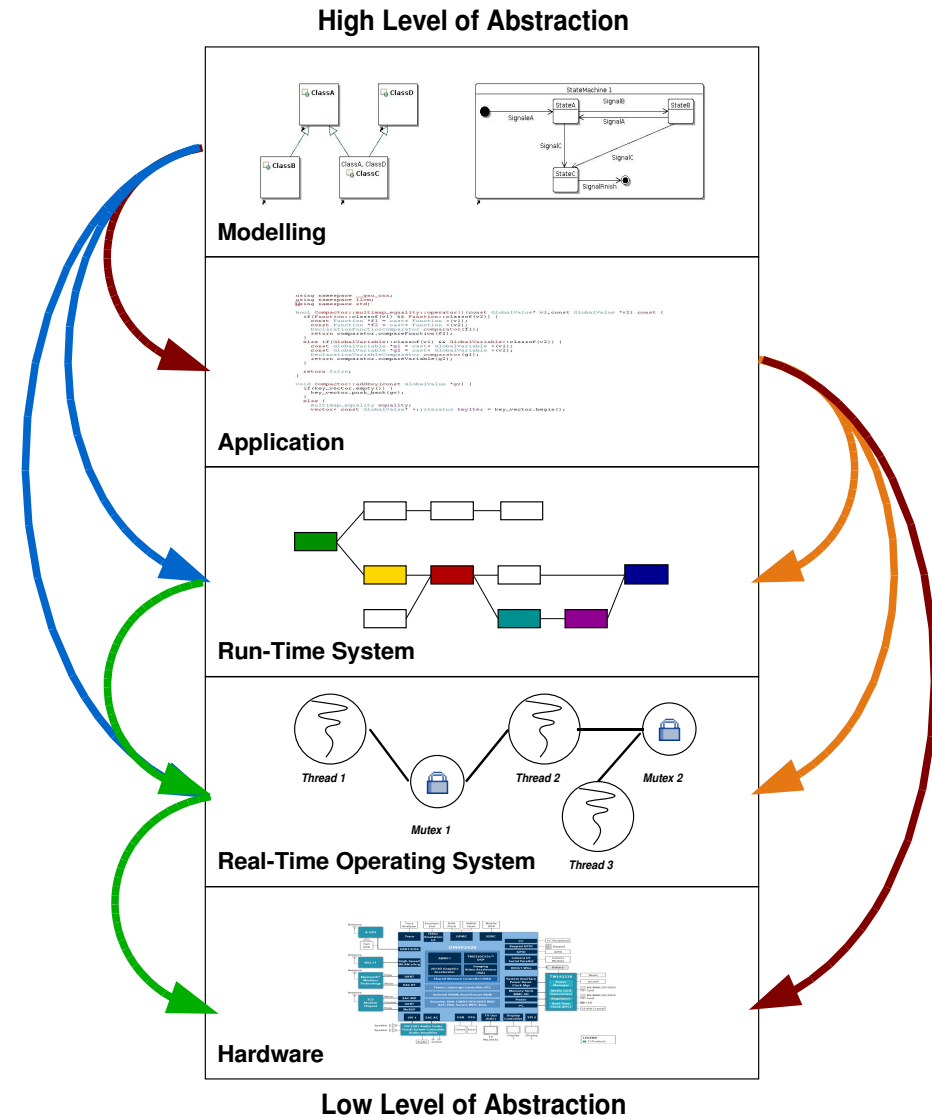
Lowering Transformations

- automated & generic
- automated
- manual



Lowering Transformations

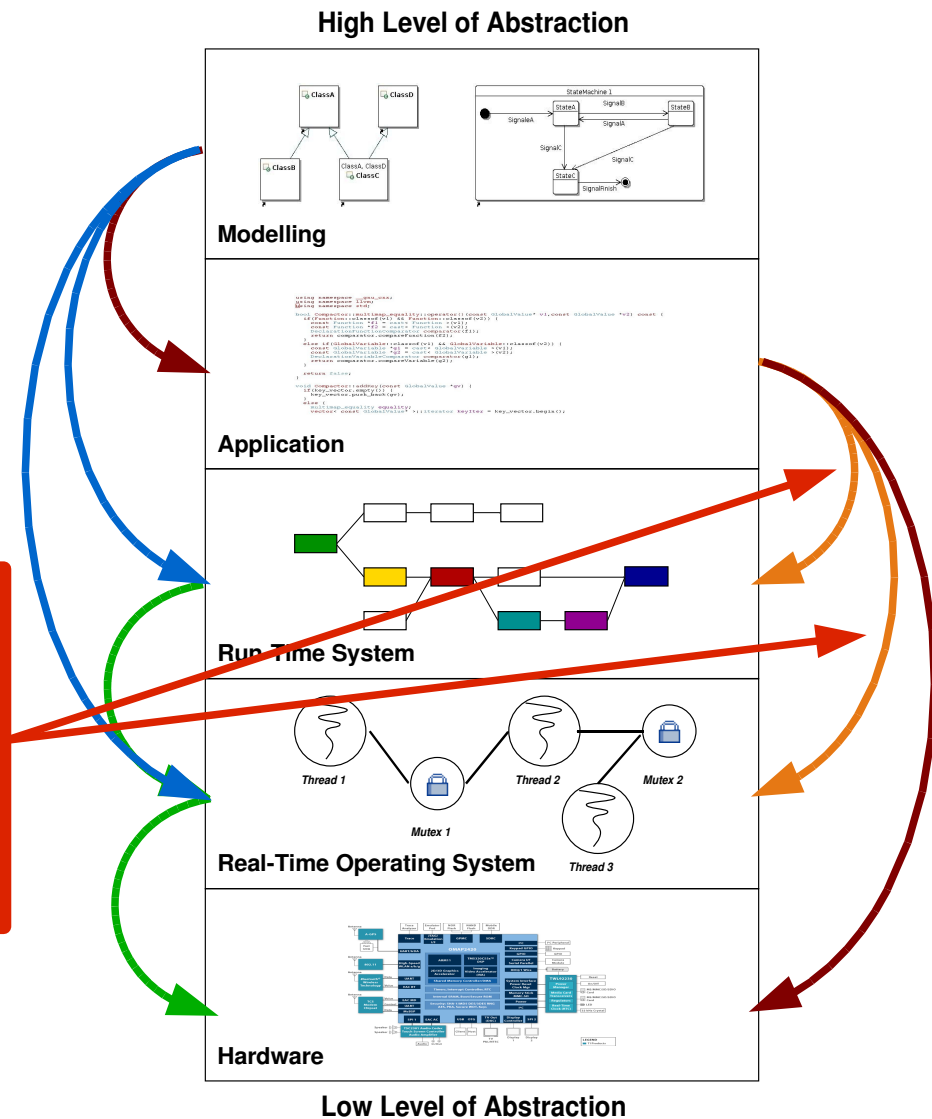
- automated & generic
- automated
- manual
- inherently manual



Lowering Transformations

- automated & generic
- automated
- manual
- inherently manual

→ automate the manual steps (not the inherently manual ones)



Overview

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Approach

- decouple Application and OS/Run-Time System
 - intermediate representation
 - independent of the employed control flow abstraction
- combination of different
 - front ends and
 - back ends
- similar to compiler construction
- intermediate representation
 - control flow graphs (CFG)
 - basic blocks

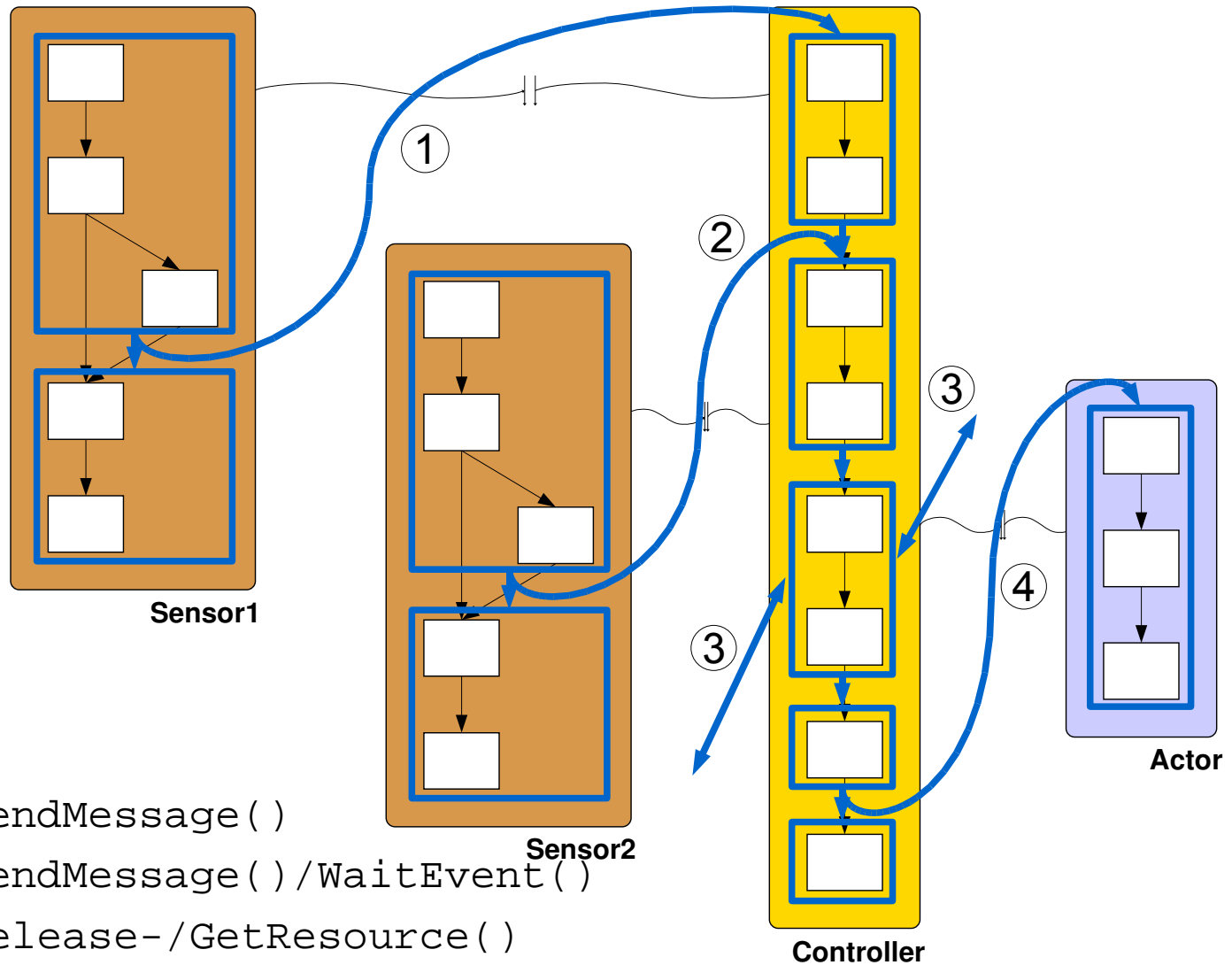


Atomic Basic Blocks

- specification of dependencies across different CFGs
 - data dependencies
 - explicitly modelled logical and temporal dependencies
 - mutual exclusion
- ABB-graph is superimposed on forest of CFGs
- ABBs aggregate several basic blocks
- ABB boundaries
 - forking/joining other CFGs
 - being joined by another CFG
 - critical sections



Atomic Basic Blocks



① `SendMessage()`

② `SendMessage()/WaitEvent()`

③ `Release-/GetResource()`

④ `SendMessage()`



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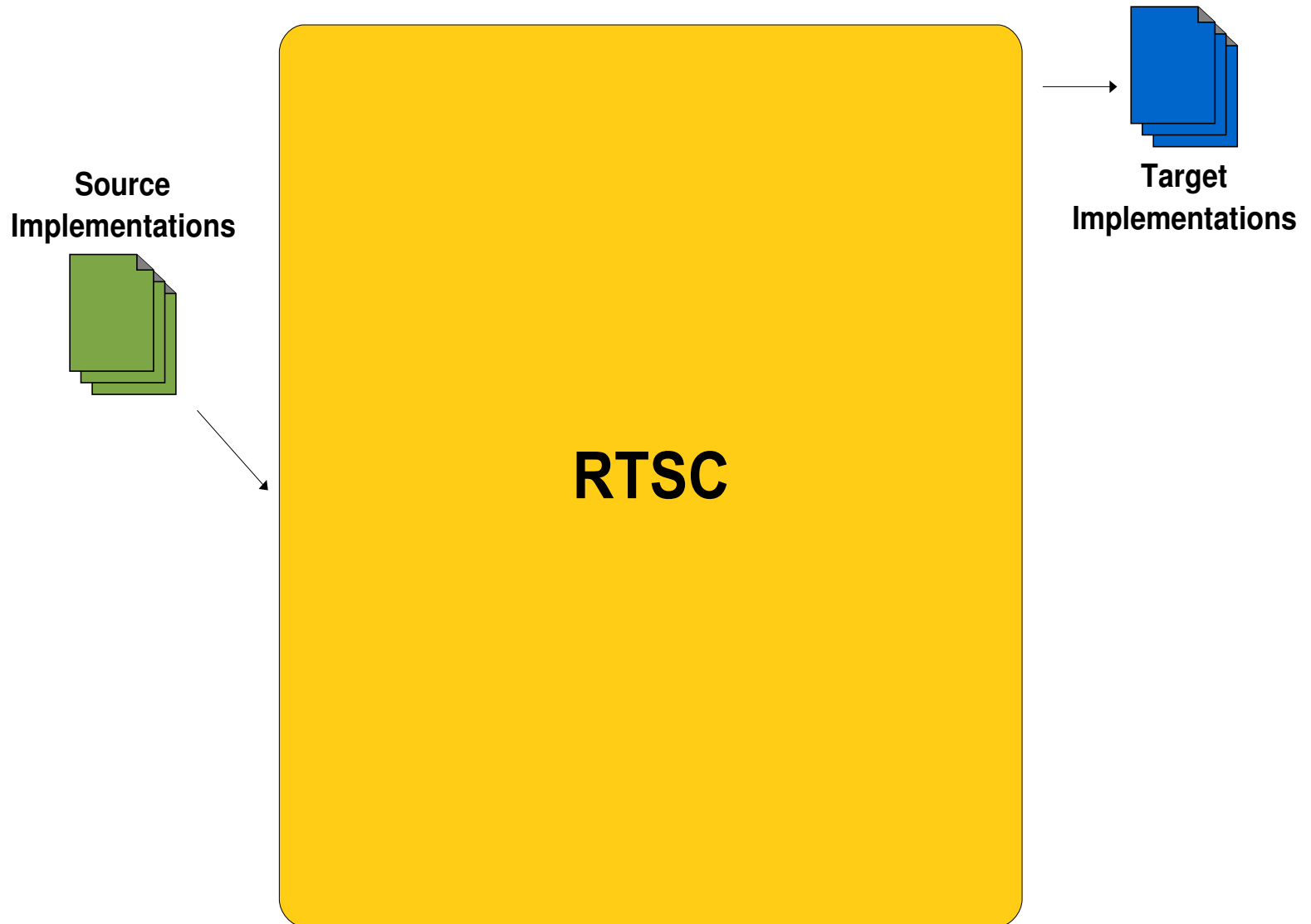


The Real-Time Systems Compiler

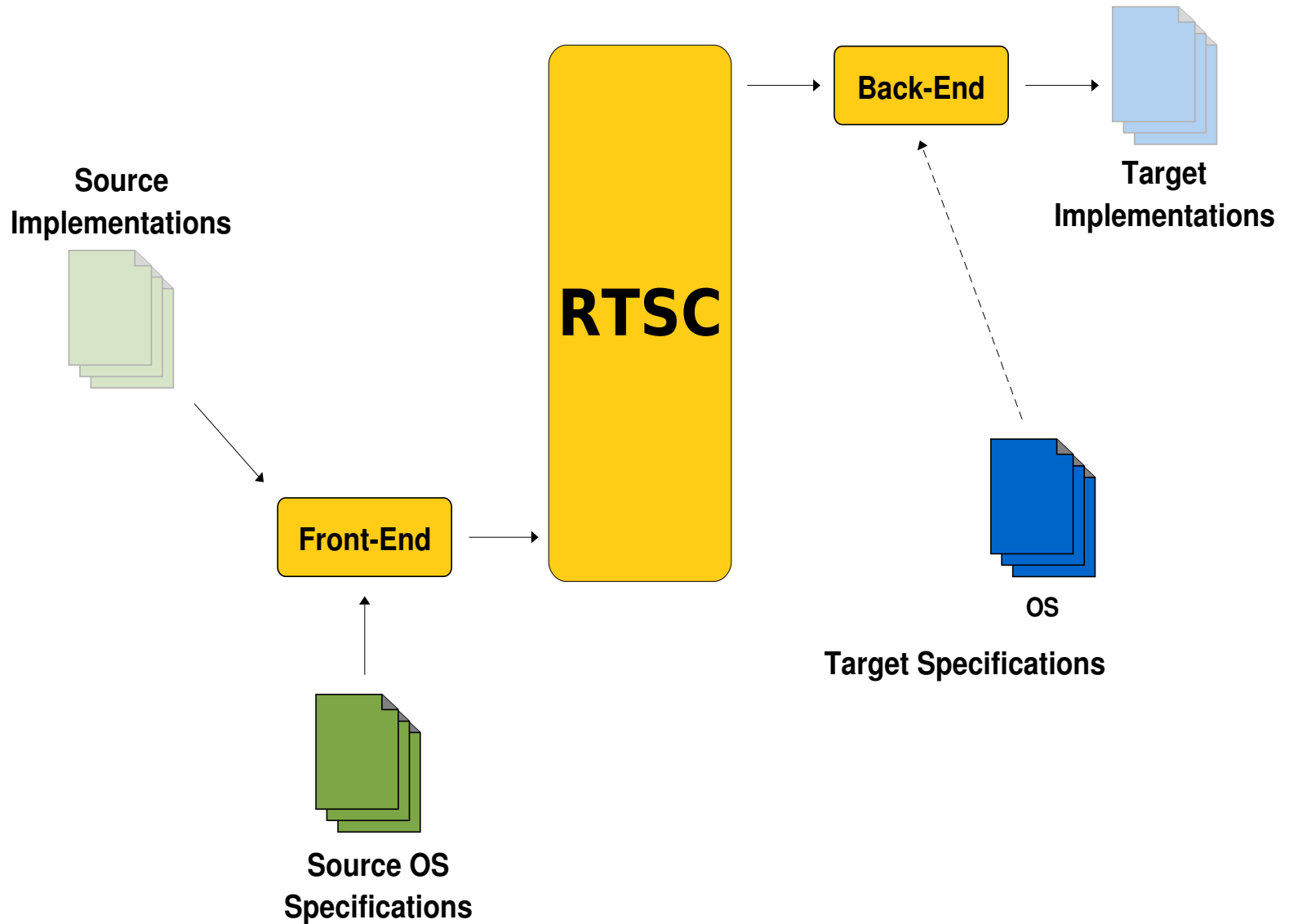
- OS aware compiler
 - based on LLVM
- uses ABBs as intermediate representation
 - ABBs are implemented on top of the LLVA
 - intermediate representation employed by the LLVM
 - typed assembler



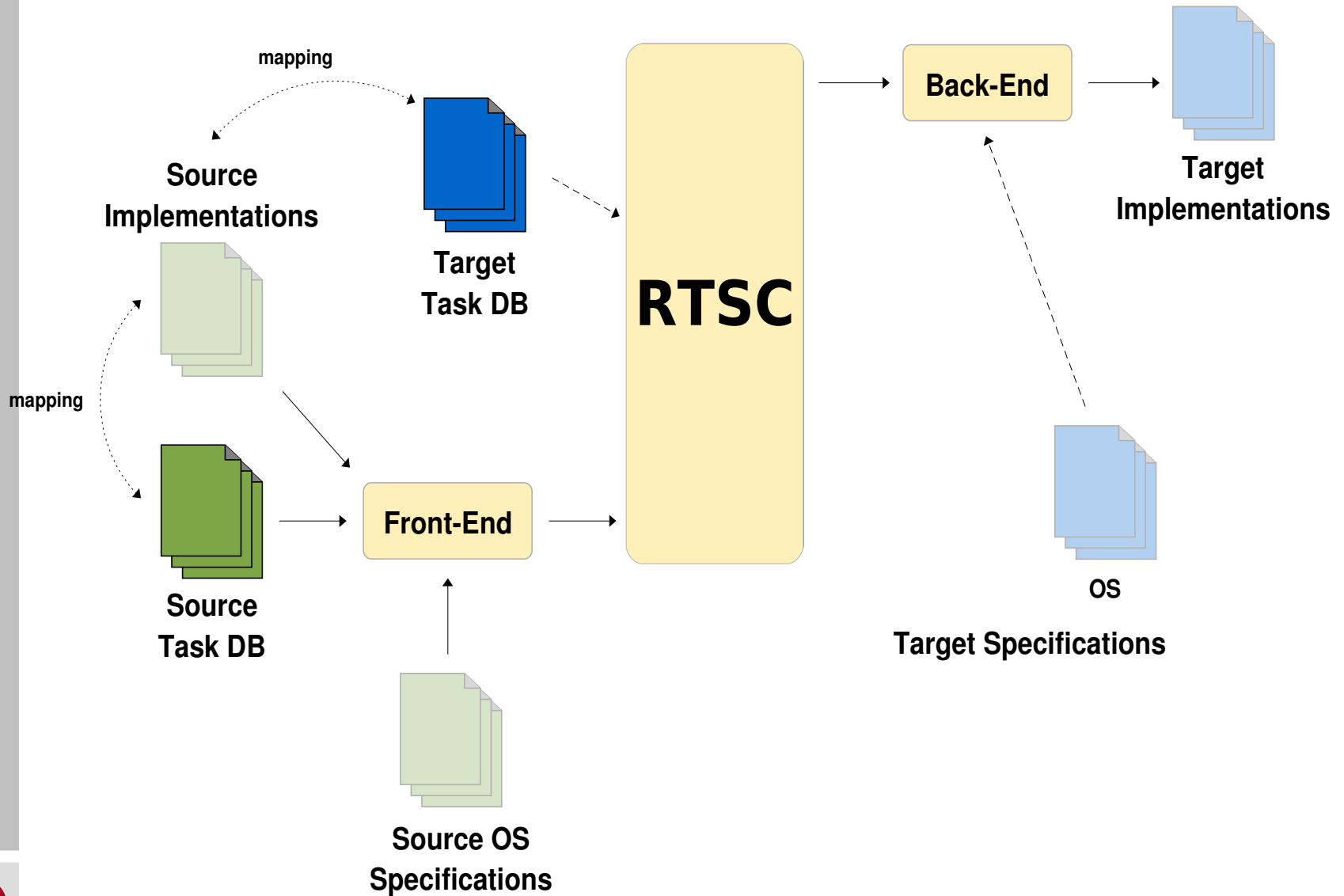
Input & Output



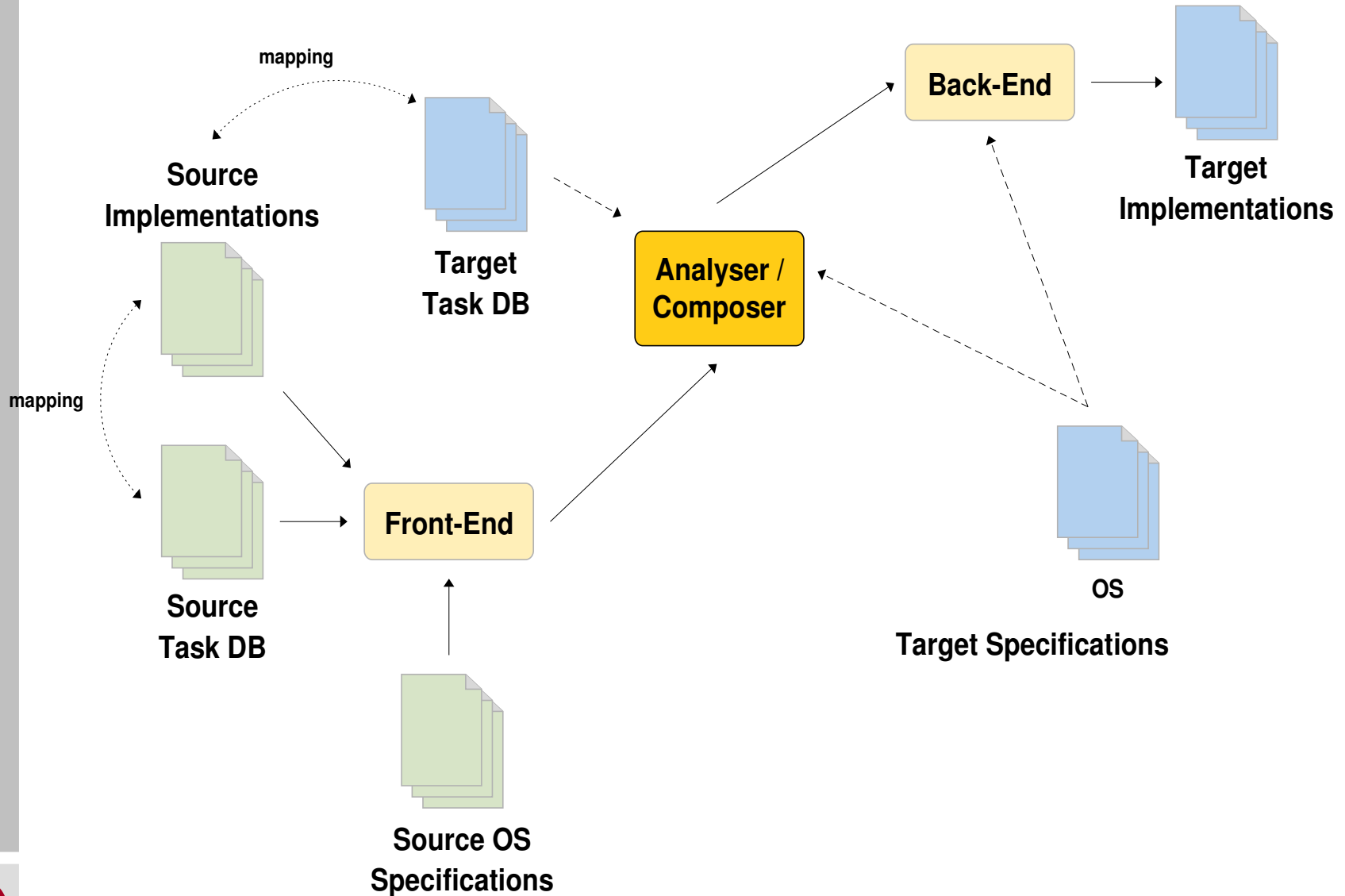
OS-specific Front- and Back-End



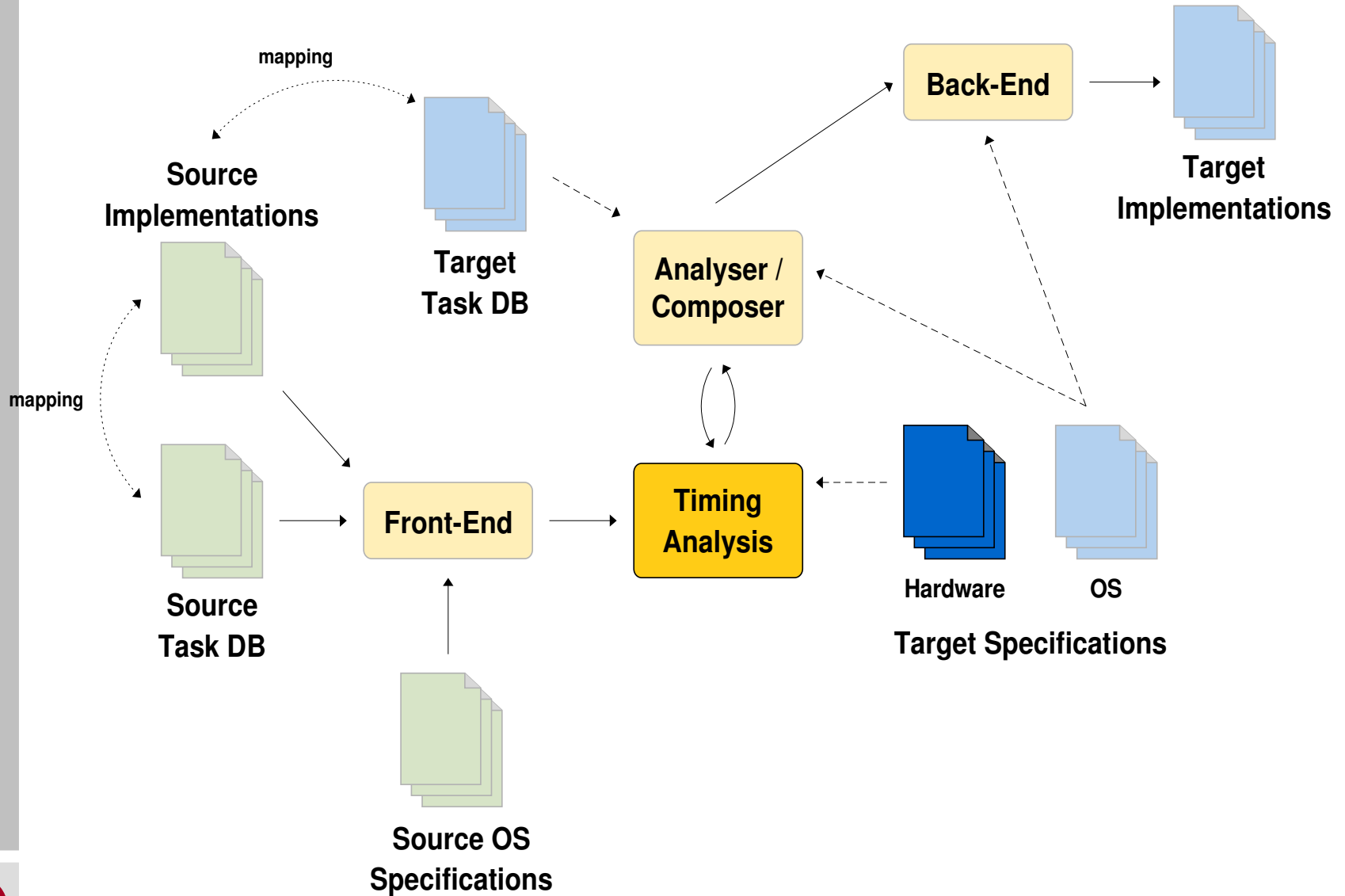
Specifying Events



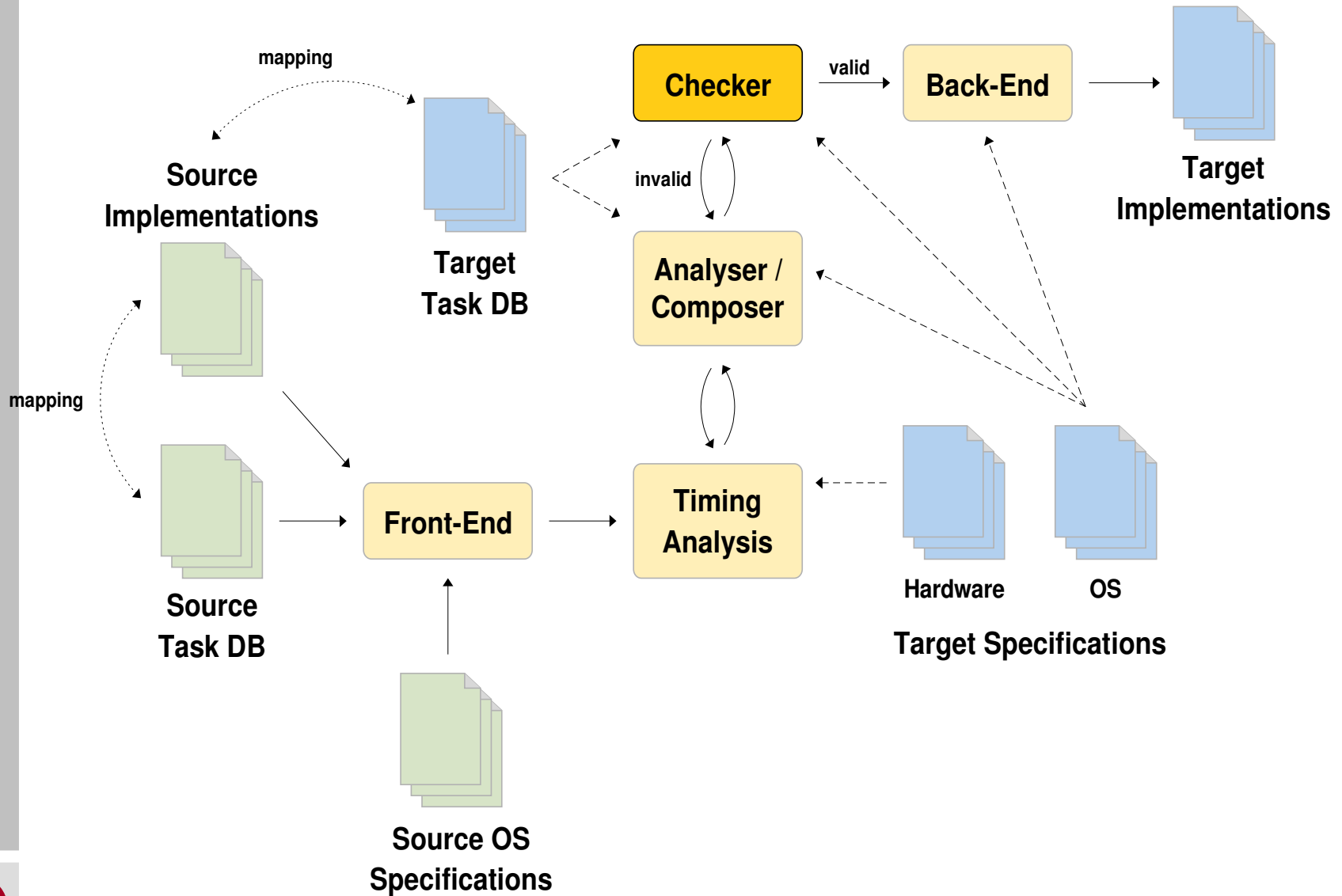
Mapping ABBs to OS mechanisms



WCET Analysis



Schedulability Analysis



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Future Work

- current status
 - early prototype of a C front-end
 - not that exciting: no ABBs except global data dependencies
 - framework for creating ABB-graphs
- future work
 - OSEK OS and OSEK ttOS front-end and back-end



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**Thank you very much
for your attention!**

